# Card-shuffling via convolutions of projections on combinatorial Hopf algebras 



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XXVII FPSAC, Daejeon, Korea
9 July 2015

## Part I: The riffle-shuffle

- Cut the deck with symmetric binomial distribution;


$$
\text { Prob }=2^{-n}\binom{n}{i}
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$$
\begin{array}{lll}
1 \oslash & 4 \uparrow & \text { Prob }=\frac{1}{4} \\
2 \diamond &
\end{array}
$$

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$$
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2 \diamond
\end{array} \quad 4 \uparrow \quad \text { Prob }=\frac{1}{3}
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$$
\text { Prob }=\frac{1}{2} \quad 1 \varnothing
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Prob $=\frac{1}{2}$

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Bayer-Diaconis (1992):

$2 \diamond$
30

Randomising $n$ distinct cards needs $\frac{3}{2} \log n$ shuffles.

## A new tool: the shuffle (Hopf) algebra $\mathcal{S}$

- graded: $\mathcal{S}=\bigoplus \mathcal{S}_{n}$
- basis of $\mathcal{S}_{n}$ is $\mathcal{B}_{n}:=\{$ words of length $n\}=\{$ decks of $n$ cards $\}$ (possibly with repeated letters / cards)


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- coproduct $\Delta: \mathcal{S} \rightarrow \mathcal{S} \otimes \mathcal{S}$ is sum of all deconcatenations

$$
\begin{aligned}
\Delta([155])= & \epsilon \otimes[155]+[1] \otimes[55]+[15] \otimes[5]+[155] \otimes \epsilon \\
& \text { empty deck }=\text { unit of } \mathcal{S}
\end{aligned}
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Relation with riffle-shuffling:
$\operatorname{Prob}(x \rightarrow y)=$ coefficient of $y$ in $\frac{1}{2^{n}} m \circ \Delta(x)$ for $x, y \in \mathcal{B}_{n}$.

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& =\frac{5}{8}[155]+\frac{2}{8}[515]+\frac{1}{8}[551]
\end{aligned}
$$

## Consequences

$$
\operatorname{Prob}(x \rightarrow y)=\text { coefficient of } y \text { in } \frac{1}{2^{n}} m \circ \Delta(x) \text { for } x, y \in \mathcal{B}_{n}
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Theorem (w/ Diaconis, Ram, 2014): Algorithm for a basis of eigenvectors of $m \circ \Delta$ on shuffle algebra, from Hopf algebraic structure theorems.

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Corollary (and folklore): Stationary distribution is uniform.
Corollary (and folklore): Start with $n$ distinct cards in ascending order. After $t$ riffle-shuffles:

$$
\text { Expect }\{\text { number of descents }\}=\left(1-\left(\frac{1}{2}\right)^{t}\right) \frac{n-1}{2} .
$$

## Other shuffling schemes

$\operatorname{Prob}(x \rightarrow y)=$ coefficient of $y$ in $\mathbf{T}(x)$ for $x, y \in \mathcal{B}_{n}$.

Riffle-shuffle

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\mathbf{T}=\frac{1}{2^{n}} m \circ \Delta
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Riffle-shuffle
Top-to-random

$$
\mathbf{T}=\frac{1}{2^{n}} m \circ \Delta
$$

$$
\mathbf{T}=\frac{1}{n} m \circ \Delta_{1, n-1}
$$

Project the coproduct to $\mathcal{S}_{1} \otimes \mathcal{S}_{n-1}$

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\Delta([155])=\emptyset \otimes[155]+[1] \otimes[55]+[15] \otimes[5]+[155] \otimes \emptyset
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Top-or-bottom-to-random $\mathbf{T}=\frac{1}{2 n}\left(m \circ \Delta_{1, n-1}+m \circ \Delta_{n-1,1}\right)$

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$$
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$$

$$
\mathbf{T}=\sum q^{i}(1-q)^{n-i} m \circ \Delta_{i, n-i}
$$

$$
(1-q)^{3} \quad q(1-q)^{2} \quad q^{2}(1-q) \quad q^{3}
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cut-and-interleave Diaconis, Fill, Pitman (1992)

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$$

descent operators
Patras (1994)

## Consequences

$\operatorname{Prob}(x \rightarrow y)=$ coefficient of $y$ in $\mathbf{T}(x)$ for $x, y \in \mathcal{B}_{n}$.

Theorem (2015): For many significant $\mathbf{T}$
(top-to-random, top-or-bottom-to-random, etc.), we can algorithmically compute an eigenbasis.

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Theorem (2015): For many significant $\mathbf{T}$ (top-to-random, top-or-bottom-to-random, etc.), we can algorithmically compute an eigenbasis.

Corollary: Stationary distribution is always uniform.
Corollary: Start with $n$ distinct cards in ascending order. After $t$ top-to-random shuffles:
Prob $\{$ descent at bottom $\}=\left(1-\left(\frac{n-2}{n}\right)^{t}\right) \frac{1}{2}$.

## Part II: Break-and-recombine other combinatorial objects

On other combinatorial Hopf algebras, define Markov chain by:
$\operatorname{Prob}(x \rightarrow y):=$ coefficient of $y$ in $\mathbf{T}(x)$ for $x, y \in \mathcal{B}_{n}$
$\mathbf{T}=$ descent operator $\sim$ style of shuffle

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\mathcal{B}_{n}=\text { basis } \sim \text { combinatorial object }
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shuffle algebra
Connes-Kreimer trees
graph Hopf algebra
symmetric functions, schur basis
$\longrightarrow$ card-shuffling
$\longrightarrow$ tree pruning
$\longrightarrow$ edge removal
$\longrightarrow$ a chain on partitions

## Example: top-to-random on partitions

$\operatorname{Prob}(\lambda \rightarrow \mu):=$ coefficient of $s_{\mu}$ in $\frac{1}{n} m \circ \Delta_{1, n-1}\left(s_{\lambda}\right)$.

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\Delta_{1, n-1}\left(s_{\lambda}\right)=\sum_{\nu=\lambda \backslash \square} s_{(1)} \otimes s_{\nu} ; \quad s_{(1)} s_{\nu}=\sum_{\mu=\nu \cup \square} s_{\mu}
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\frac{1}{3} m \circ \Delta_{1,2}(\square \square \square)=\frac{1}{3} m(\square \otimes \square \square)
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|  |  | $\mu$ |  |  |
| :---: | :---: | :---: | :---: | :---: |
|  |  | $(3)$ | $(2,1)$ | $(1,1,1)$ |
| $\lambda$ | $(3)$ | $1 / 3$ | $1 / 3$ |  |
|  | $(2,1)$ | $1 / 3$ | $2 / 3$ | $1 / 3$ |
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## Example: top-to-random on partitions

## $\operatorname{Prob}(\lambda \rightarrow \mu):=$ coefficient of $s_{\mu}$ in $\frac{1}{n} m \circ \Delta_{1, n-1}(S \lambda)$.

Divide $s_{\lambda}$ by number of standard tableaux of shape $\lambda$

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To make coefficients sum to 1 , use "the Doob transform"

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|  |  | $(3)$ | $(2,1)$ | $(1,1,1)$ |
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|  | $(2,1)$ | $1 / 6$ | $2 / 3$ | $1 / 6$ |
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## General theorems

$$
\operatorname{Prob}(x \rightarrow y):=\text { coefficient of } \frac{y}{\eta(y)} \text { in } \mathbf{T}\left(\frac{x}{\eta(x)}\right) \text { for } x, y \in \mathcal{B}_{n}
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$\frac{j}{n}$, \#partitions with $j$ parts of size 1

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Quotient / sub Hopf algebras give lumpings (independent of $\mathbf{T}$ ) RSK shape of top-to-random-with-standardisation = top-to-random on partitions, because $\Lambda \longleftrightarrow$ FSym $\hookrightarrow$ FQSym

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A Please tell me your favourite Hopf algebras and non-negative linear maps
pf size 1

## Thank you!

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